#### Logical Languages part 2 2020

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#### Formal logic and intro to predicate calculus

All propositions can be expressed in clausal form

$$B_1 \cup B_2 \cup \ldots \cup B_n \subset A_1 \cap A_2 \cap \ldots \cap A_m$$

Right side implies left side If all of the A are true, at least one B is true

#### Predicate calculus and proving theorems

 One way to simplify resolution process: restrict to simpler forms of propositions

#### **Horn Clause**

Either (1) single atomic proposition on left side (2) empty left side

Also called (1) Headed horn clause (2) Headless Horn clause

#### Predicate calculus and proving theorems

 One way to simplify resolution process: restrict to simpler forms of propositions

#### **Horn Clause example:**

(1) Headed horn clause

likes (bob, trout)  $\subset$  likes (bob, fish)  $\cap$  fish (trout)

(2) Headless Horn clause

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father (bob, jake) Often used to state fact

Most, but not all, propositions can be stated as Horn clauses. The restriction to Horn clauses makes resolution a practical process for proving theorems.

#### Predicate calculus and proving theorems

- Main idea: Presence of variables in propositions requires resolution to find values for the variables that allows matching to succeed
- Unification: Finding values for variables in propositions that allows matching to succeed
- Instantiation: temporary assigning of values to variables to allow unification
- Backtracking: if resolution process to instantiate a variable with a value fails to complete required matching, then we backtrack and instantiate variable with different value

We will use:

http://www.swi-prolog.org/

https://www.swi-prolog.org/pldoc/doc\_for?object=manual

You can run Prolog programs in several ways:

If you want to install it in your own machine, please install the latest version of the SWI-prolog. In this case please download it from http://www.swiprolog.org/

Our lab also has the SWI-prolog. You can use it from there. To use it, first you need to login into the lab machines. Now you are ready to work with Prolog.

Lets create a sample Prolog program file name simple.pl which has following two lines

person(bob).

father(bob,sam).

Now to open the Prolog complier write command swipl.

Now to load the simple.pl file type ['simple.pl']. in the compiler terminal. The "." at the end of ['simple.pl']. is used to mark the end of the command in Prolog.

Then you can query based on this simple file: person(bob).

returns true

father(bob,X).

returns X = sam.

There are other instructions available:

http://www.cs.toronto.edu/~sheila/324/f05/tuts/swi.pdf Ctrl-d to guit

- University of Aix-Marselle (NLP) and Edinburgh (automated theorem proving) in mid 1970s
- Prolog dialect has several forms. Here we focus on Edinburgh syntax

Terms: 1. constant, 2. variable, 3. structure

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- String of letters, digits, underscores beginning with lower case letter Examples? likes, father, my\_classes
- String of printable ASCII characters Examples? , :- have predefined meanings

- Terms: 2. variable
- String of letters, digits, and underscores, beginning with Uppercase letters

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Examples: X, List

- Terms: 2. variable
- String of letters, digits, and underscores, beginning with Uppercase letters
- Variables are not bound to types by declaration

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- Variable instantiation: Binding a variable to a value and thus to a type

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- Variable instantiation: Binding a variable to a value and thus to a type
- Lasts only as long as it takes to satisfy one complete goal (proof or disproof of proposition)
- Example: student(X) instantiation will set a variable X to bob and check the proposition student(bob)

- Terms: 2. variable
- Prolog variables only **distant relatives** to imperative languages both in semantics and use

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**Used for?** 

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- Terms: 3. structure
- Represents atomic propositions of predicate calculus
  - functor(parameter\_list)
  - Example? father(jon, shelley)
  - used to specify facts in Prolog
  - > also a predicate when specifying a question (query)

Fact statements

 Construct hypotheses or database of assumed information; statements from which new information can be inferred

#### Fact statements

- Construct hypotheses or database of assumed information; statements from which new information can be inferred
- Remember: facts we have in database; then queries/goals asking about database

Statement forms

1. Headless Horn clauses of predicate calculus

Examples??

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female(shelley)
male(bill)
father(bill, jake)

Statement forms

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Examples??

female(shelley)
male(bill)
father(bill, jake)

Why are the first letters of each term lower case?

Statement forms

1. Headless Horn clauses of predicate calculus

Examples??

female(**s**helley) male(**b**ill) father(**b**ill, **j**ake)

Why are the first letters of each term lower case? Answer: these are not variables, but facts (or queries)

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Statement forms

2. Rule statements (these will correspond to headed horn clauses)
# Prolog

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consequence :- expression

Here the expression implies consequence (**right side implies left side**)

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consequence :- expression

Here the expression implies consequence (right side implies left side)

The expression can be a single term or conjunction Example: female(shelley), child(shelley)

# Prolog

Statement forms

2. Rule statements (these will correspond to headed horn clauses)

Example headed horn clauses:

ancestor(mary, shelley) :- mother(mary, shelley)

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Example headed horn clauses:

ancestor(mary, shelley) :- mother(mary, shelley)

Reads: If mary is the mother of shelley, then this implies that mary is an ancestor of shelley

2. Rule statements (these will correspond to headed horn clauses)

# **Use of Variables in Prolog statements**

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# Use of Variables in Prolog statements parent(X,Y) :- mother(X,Y)

Meaning?

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#### **Use of Variables in Prolog statements**

parent(X,Y) :- mother(X,Y)

Meaning? If there are instantiations of X, Y such that mother(X,Y) is true, then for those instantiations of X and Y, parent (X,Y) is true

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#### **Use of Variables in Prolog statements**

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Meaning? If there are instantiations of X, Y such that mother(X,Y) is true, then for those instantiations of X and Y, parent (X,Y) is true

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There could be several X,Y pairs in the database for which parent(X,Y) is true. jon, shelley mary,liz etc

2. Rule statements (these will correspond to headed horn clauses)

### **Use of Variables in Prolog statements**

parent(X,Y) :- mother(X,Y)

Meaning? If there are instantiations of X, Y such that mother(X,Y) is true, then for those instantiations of X and Y, parent (X,Y) is true

Use of variables allows to generalize meanings

3. Goal statements (these will correspond to headless horn clauses, like the fact statements)

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3. Goal statements (these will correspond to headless horn clauses)

- So far: we have described statements as logical propositions, for facts and logical relationships between facts. These are the basis for theorem proving.
- The theorem: in the form of a proposition that we want to prove or disprove (called goals or queries) Example: man(fred)

3. Goal statements (these will correspond to headless horn clauses)

Example: man(fred)

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The system will respond either:

true: proved goal and true under database of facts and relations

false: either goal was determined as false, or system was unable to prove it

# Prolog

Statement forms

3. Goal statements (these will correspond to headless horn clauses)

Another example: father(X, mike)

# Prolog

Statement forms

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Another example: father(**X**, mike)

Note that **X** is a variable (starts with capital letter)

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Another example: father(**X**, mike)

Note that **X** is a variable (starts with capital letter)

When a variable is present, the system not only asserts validity, but **identifies instantiations of variable that make goal true** 

```
simple.pl file includes:
```

```
% Simple example for testing
% swipl from command line
% Inside compiler:
% ['simple.pl'].
% person(bob).
% returns true
% father(bob,X).
% returns X = sam.
% control d to exit
```

person(bob).
father(bob,sam).

simple.pl let's try it in compiler:

swipl from command line

Inside compiler: ['simple.pl'].

Notice we always have a period after statement

simple.pl let's try it in compiler:

- Inside compiler: ['simple.pl'].
- > person(bob).
   Returns?

simple.pl let's try it in compiler:

- Inside compiler: ['simple.pl'].
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   Returns?

simple.pl let's try it in compiler:

- > Inside compiler: ['simple.pl'].
- > father(bob,X).
   Returns? X = sam.

simplemore.pl let's add more facts to file:

person(bob). father(bob,sam). father(sam,liz).

simplemore.pl let's add more facts to file:

```
person(bob).
father(bob,sam).
father(bob,liz).
```

```
> father(bob,X).
    Returns?
```

initially returns X = sam

simplemore.pl let's add more facts to file:

```
person(bob).
father(bob,sam).
father(bob,liz).
```

```
> father(bob,X).
    Returns?
```

initially returns X = sam
Type ; and will return next item here:
X = liz

```
Prolog demos
```

```
simplemore.pl let's add more facts to file:
```

```
person(bob).
father(bob,sam).
father(bob,liz).
```

```
> father(bob,X).
    Returns?
```

returns X = sam ; X = liz

So system will attempt (called unification) to find instantiations of X that results in true value for goal

### simple2.pl

%http://faculty.otterbein.edu/psanderson/csc326/notes/Prolo gNotes.html

```
mother(iva, pete).
mother(iva, ed).
mother(iva, becky).
mother(kay, nancy).
mother(kay, bob).
mother(kay, diane).
mother(becky, katie).
husband(dwight, iva).
husband(robert, kay).
husband(pete, nancy).
```

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husband(pete, nancy).
```

Things to try: mother(kay, nancy). mother(kay, kay). mother(kay, Who). press ;

# simple2.pl

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%http://faculty.otterbein.edu/psanderson/csc326/notes/Prolo gNotes.html

```
mother(iva, pete).
mother(iva, ed).
mother(iva, becky).
mother(kay, nancy).
mother(kay, bob).
mother(kay, diane).
mother(becky, katie).
husband(dwight, iva).
husband(robert, kay).
husband(pete, nancy).
```

```
Things to try:
mother(kay, Who). press ;
mother(kay,Who).
Who = nancy ;
Who = bob ;
Who = diane.
```

#### Note about form

 Goal and non goal statements (e.g., facts, rules) can have the same form

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- So Prolog implementation must have means to differentiate goals and non goals
- We separated by reading in facts file first

### Inferencing process of Prolog

 Prolog resolution is critical (proving true, or false cannot prove) Inferencing process of Prolog

- Prolog resolution is critical (proving true, or false cannot prove)
- Queries are called goals
   If a goal is a compound proposition, it consists of subgoals
- To prove goal true: inferencing process must find chain of rules and/or facts in the database

Inferencing process of Prolog

If Q is a goal, then either Q must be found in the database, or inferencing must find fact P1 and propositions P2, P3, P4, ... Pn such that:

Q :- Pn

. . .
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If Q is a goal, then either Q must be found in the database, or inferencing must find fact P1 and propositions P2, P3, P4, ... Pn such that:



Process is called matching, satisfying, or resolution

man(bob) query

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Easy if database includes this fact; then proof trivial

```
Inferencing process of Prolog. Example:
```

```
man(bob) query
```

More complex if database includes rules:

```
father(bob).
man(X) :- father(X).
```

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man(bob) query
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```

Prolog needs to find the two statements, and use them to infer the truth of the goal; needs unification to instantiate X temporarily to bob

```
Inferencing process of Prolog. Example:
```

```
man(X) query
```

More complex if database includes rules:

```
father(bob).
father(jon).
man(X) :- father(X).
```

Prolog must match goal against propositions in database

e.g., first find bob then jon (remember we used ; in compiler)

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man(bob) query

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```
Database includes rules:
father(bob).
man(X) :- father(X).
```

How does Prolog do it? Two possibilities:

**1. Forward chaining: search for and find first proposition** father(bob); goal is inferred by matching first proposition with right side of second rule father(X) through instantiation of X to bob, and then matching left side of second proposition to goal man(bob)

man(bob) query

Database includes rules: father(bob). man(X) :- father(X).

How does Prolog do it? Two possibilities:

**2. Backward chaining: first match goal** with left side of second proposition man(X) through the instantiation of X to bob; as last step, match right side of second proposition (now father(bob)) with first proposition

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**1. Forward chaining:** search for and find first proposition father(bob); goal is inferred by matching first proposition with right side of second rule father(X) through instantiation of X to bob, and then matching left side of second proposition to goal man(bob)

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Which does Prolog use?

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1. Forward chaining: search for and find first proposition father(bob); goal is inferred by matching first proposition with right side of second rule father(X) through instantiation of X to bob, and then matching left side of second proposition to goal man(bob)

**2. Backward chaining:** first match goal with left side of second proposition man(X) through the instantiation of X to bob; as last step, match right side of second proposition (now father(bob)) with first proposition

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Prolog uses Backward chaining. First match goal.