# Programming Languages Scheme part 4 and other functional 

 2020
## Lots of equalities!

## Summary:

- eq? for symbolic atoms, not numeric (eq? 'a 'b)
- $=$ for numeric, not symbolic (=57)
- eqv? for numeric and symbolic


## equal versus equalsimp

(define (equalsimp lis1 lis2) (cond
((null? lis1) (null? lis2))
((null? lis2) \#f)
((eq? (car lis1) (car lis2))
(equalsimp (cdr lis1)
(cdr lis2)))
(else \#f)
(define (equal lis1 lis2) (cond
((not (list? lis1)) (eq? lis1
lis2))
((not (list? lis2)) \#f)
((null? lis1) (null? lis2))
((null? lis2) \#f)
((equal (car lis1) (car
lis2))
(equal (cdr lis1) (cdr
lis2)))
(else \#f)
)
)

## append

## (define (append lis1 lis2) <br> (cond <br> ((null? lis1) lis2) <br> (else (cons (car lis1) (append (cdr lis1) lis2))) <br> ))

- Reminding ourselves of cons (run it on csi):
(cons '(a b) '(c d)) returns ((a b) c d)
(cons '((ab) c) '(d (e f))) returns (((ab)c)d(ef))


## Adding a list of numbers

- This works: (+ 3710 2)
- This doesn't work: (+ (3 710 2))
- Why?


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How would we achieve the second option?
Breakout groups

## Adding a list of numbers

- We want: (+ (3 710 2))

```
(define (adder a_list)
    (cond
((null? a_list) 0)
(else (eval(cons '+ a_list)))
)
)
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We'll do a little "trick" ...

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- Why the quote on ' + ?


## Adding a list of numbers

- We want: (+ (3 710 2))
(define (adder a_list)
(cond
((null? a_list) 0) (else (eval(cons '+ a_list)))
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- cons creates new list with + and a_list
- Why the quote on ' + ?
- Quote so that eval will not evaluate in evaluation of cons


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(define (adder a_list)
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```

- Adder (+ 123 4)
- Calls (eval (+ 123 4))
- And returns (+ 123 4)


## Adding a list of numbers

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(define (adder a_list)
    (cond
    ((null? a_list) 0)
    (else (eval(cons '+ a_list)))
    )
)
```

- Create adder function and load into csi
- Run on sci adder (+ 123 4)
- Run on sci (eval (+ 123 4))


## Adding a list of numbers

- We want: (+ (3 710 2))

```
(define (adder a_list)
    (cond
((null? a_list) 0)
(else (eval(cons '+ a_list)))
)
)
```

Examples:
(adder '(1 2 3))

## Adding a list of numbers

- We want: (+ (3 710 2))

Let's each write another way of doing this...

Create adder2 function and load into csi
Run on sci (adder2 '(3 710 2))

## Adding a list of numbers

- We want: (+ (3 710 2))

Let's each write another way of doing this... Hint: use car and cdr

Create adder2 function and load into csi
Run on sci (adder2 '(3 710 2))

Other functional languages

## Common LISP

- Combination of many features of popular dialects of LISP, early 1980s
- Large and complex language, opposite of Scheme
- Features include: records; arrays; complex numbers; character strings; iterative control statements; etc.
- So not purely functional, has imperative features


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## ML Language

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## What were those?

## ML Language

- Syntax closer to Pascal and other imperative than to LISP
- Strongly typed (whereas Scheme is essentially typesless) with no type coercions
- Has identifiers, but once set cannot be changed more like final declarations in Java or const in C/C++


## Functional declarations ML

- Format:
fun name (parameters) = body;


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Example (run it):
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## Example:

fun circumf(r) $=3.14^{*} r^{*} r$;
The type here is inferred as float from the type of the literal in the expression

## Functional declarations ML

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Example:
fun times10(x) = 10*x;
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Also inferred as int (default type)
What happens if called with square(2.75)???

## ML Language

https://www.tutorialspoint.com/execute_smInj_online.php

## ML Language

- Try running some code:
fun times10(x) = 10*x; times10(5);


## ML Language

- Try running some code:
fun times10(x) $=10 * x$; times10(5);
times10(5.1);
What happens???


## ML Language

- Try running some code:
fun times10(x) = 10*x; times10(5);
times10(5.1);
Yields error; expecting int...
It's strongly typed!!


## ML Language

- We could also specify type.
fun square( $x$ :real $)=x * x$;


## ML Language

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fun square( $x$ :real $)=x * x$;
Enough to infer that type is real


## ML Language

- These are all valid:
fun square( $x:$ real $)=x * x$;
fun square(x) = (x:real) * x;
fun square $(x)=x^{*}$ (x:real)
Enough to infer that type is real


## ML Language

- These are all valid:
fun square(x:real) $=x$ * $x$;
fun square(x) = (x:real) * x;
fun square $(x)=x^{*}$ (x:real)
Enough to infer that type is real
Type inference also used in Haskell, Miranda, F\#


## ML Language

- Try running some more code:
fun square(y:real) $=y^{*} y$;
square(5.1);
square(5.0);


## ML Language

- What about this?
fun square(y:real) $=y^{*} y$; square(5);


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fun square(y:real) $=y^{*} y$; square(5);

Oops another type error...

## ML Language

- What about this?
fun square(y:real) $=y^{*} y$; square(5);

Oops another type error...
Note: user defined overloaded functions not allowed, so if we wanted a square function, one for real and one for int, would have to use different names...

## ML selection

- if else format:
if expression then expression
else else_expression


## ML selection

- Example:
fun fact (n:int) $=$
if $n<=1$ then 1
else $\mathrm{n} *$ fact( $\mathrm{n}-1$ );


## ML selection

- Example:
fun fact ( $n$ :int) $=$
if $n<=1$ then 1
else $\mathrm{n} *$ fact( $\mathrm{n}-1$ );
fact(4);
Run it...


## ML selection

- Another way: pattern matching! fun fact(0) $=1$


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fun fact(0) $=1$
$\mid \operatorname{fact}(1)=1$


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fun fact(0) $=1$
$\mid \operatorname{fact}(1)=1$
| fact(n:int) $=n *$ fact( $n-1)$;


## ML selection

- Another way: pattern matching!

$$
\text { fun fact }(0)=1
$$

$\mid \operatorname{fact}(1)=1$
| fact(n:int) $=n *$ fact( $n-1$ );
Meant to mimic conditional functional definitions in math...

## ML selection

- Another way: pattern matching!
fun $\operatorname{fact}(0)=1$
$\mid \operatorname{fact}(1)=1$
| fact(n:int) $=n *$ fact( $n-1$ );
Meant to mimic conditional functional definitions in math...

If param is int that is not 0 or 1 then third definition is used...

## ML selection

- Another way: pattern matching!
fun $\operatorname{fact}(0)=1$
$\mid \operatorname{fact}(1)=1$
| fact(n:int) = n*fact(n-1);
Note that don't need the int here since it is the default


## ML selection

- Another way: pattern matching!
fun fact(0) $=1$
$\mid \operatorname{fact}(1)=1$
$\mid \operatorname{fact}(\mathrm{n})=\mathrm{n} * \operatorname{fact}(\mathrm{n}-1)$;
So this is also OK


## ML selection

- Another way: pattern matching!
fun fact(0) $=1$
$\mid \operatorname{fact}(1)=1$
$\mid \operatorname{fact}(\mathrm{n})=\mathrm{n} * \operatorname{fact}(\mathrm{n}-1)$;
fact(4)
Let's try running code above...


## ML list operations

- hd, tl are ML's version of Scheme CAR, CDR


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- Literal lists in brackets [3,5,7]; [] empty list
- :: used for cons
$4::[3,5,7]$ evaluates to?
[4,3,5,7]


## ML list operations

- Try running these:

4::[3,5,7]
hd([4,3,5,7])
$\operatorname{tl}([4,3,5,7])$

## ML list operations

- Number of elements in a list
fun length([]) $=0$


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- Number of elements in a list
fun length([]) $=0$
| length(h::t) = 1 + length(t);


## ML list operations

- Number of elements in a list
fun length([]) $=0$
| length(h::t) = 1 + length(t);
length([1,3,5])
Try running it


## ML list operations

- Append function
fun append ([],lis2) = ?
(what should we write here?)


## ML list operations

- Append function
fun append ([],lis2) $=$ lis2


## ML list operations

- Append function
fun append ([],lis2) $=$ lis2
| append(h::t,lis2) = ?
What should we do?


## ML list operations

- Append function
fun append ([],lis2) $=$ lis2
| append(h::t,lis2) = h::?


## ML list operations

- Append function
fun append ([],lis2) $=$ lis2
| append(h::t,lis2) = h::append(t,lis2);


## ML list operations

- Append function
fun append ([],lis2) $=$ lis2
| append(h::t,lis2) = h::append(t,lis2);
append([1,2],[3,4]);
Try running it...


## Let's remind ourselves Scheme

## (define (append lis1 lis2) <br> (cond <br> ((null? lis1) lis2) <br> (else (cons (car lis1) (append (cdr lis1) lis2))) <br> ))

- Reminding ourselves of cons (run it on csi):
(cons '(a b) '(c d)) returns ((a b) c d)
(cons '((ab) c) '(d (e f))) returns (((ab)c)d(ef))


## ML versus Scheme append

> fun append $([]$, lis2 $)=$ lis2
> | append(h::t,lis2) $=$
> h::append(t,lis2);
(define (append lis1 lis2)
(cond
((null? lis1) lis2)
(else (cons (car lis1)
(append (cdr lis1) lis2)))
))

## ML list operations

- Let's each try fun adder adder([1,2,3]) should return 6


## ML list operations

- Let's each try fun adder
fun adder([]) $=0$
| adder (h::t)=h+adder(t);
adder([1,2,3,4,5]);


## Names bound to values (constants)

- Format:
val new_name = expression;


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- Format:
val new_name = expression;
Binds the value to name once and cannot be rebound (nothing like an assignment statement in an imperative language!)


## Names bound to values (constants)

- Format:
val new_name = expression;
Example: usually used with a let statement:
fun area(radius) =
let val radius $=2.7$
val pi $=3.14159$
in pi*radius*radius
end;


## Higher order functions

- map
$\operatorname{map}\left(f n x=>x^{*} x^{*} x\right)[1,3,5]$;


## Higher order functions

- map
$\operatorname{map}\left(f n x=>x^{*} x^{*} x\right)[1,3,5] ;$
Note: different interpreters have slightly different notation; book notation different


## Higher order functions

- Composing two functions
$h=f o g$
(lower case o)


## Higher order functions

- Composing two functions
$h=f o g$
Example: (run it)
fun times10(x) = 10*x; times10(5);
fun plus3(y) $=3+y$;
plus3(4);
val h = times10 o plus3; h(7)

