Programming Languages Scheme part 4 and other functional 2020

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Lots of equalities!

Summary:

- eq? for symbolic atoms, not numeric (eq? `a `b)
- for numeric, not symbolic (= 5 7)
- eqv? for numeric and symbolic

```
(define (equalsimp lis1 lis2)
  (cond
      ((null? lis1) (null? lis2))
      ((null? lis2) #f)
      ((eq? (car lis1) (car lis2))
            (equalsimp (cdr lis1)
      (cdr lis2)))
      (else #f)
    )
)
```

(define (equal lis1 lis2) (cond ((not (list? lis1)) (eq? lis1 lis2)) ((not (list? lis2)) #f) ((null? lis1) (null? lis2)) ((null? lis2) #f) ((equal (car lis1) (car lis2)) (equal (cdr lis1) (cdr lis2))) (else #f)

append

```
(define (append lis1 lis2)
(cond
  ((null? lis1) lis2)
    (else (cons (car lis1)
        (append (cdr lis1) lis2)))
))
```

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Reminding ourselves of cons (run it on csi):

(cons `(a b) `(c d)) returns ((a b) c d)

(cons `((a b) c) `(d (e f))) returns (((a b) c) d (e f))

- This works: (+ 3 7 10 2)
- This doesn't work: (+ (3 7 10 2))
- Why?

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- This doesn't work: (+ (3 7 10 2))

How would we achieve the second option?

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- This doesn't work: (+ (3 7 10 2))

How would we achieve the second option? Breakout groups

```
(define (adder a_list)
  (cond
  ((null? a_list) 0)
  (else (eval(cons '+ a_list)))
)
```

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```
We'll do a little "trick" ...
```

We want: (+ (3 7 10 2))

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- Why the quote on `+?

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```

- cons creates new list with + and a_list
- Why the quote on `+?
- Quote so that eval will not evaluate in evaluation of cons

```
(define (adder a_list)
  (cond
  ((null? a_list) 0)
  (else (eval(cons '+ a_list)))
)
```

- Adder (+ 1 2 3 4)
- Calls (eval (+ 1 2 3 4))
- And returns (+ 1 2 3 4)

```
(define (adder a_list)
  (cond
  ((null? a_list) 0)
  (else (eval(cons '+ a_list)))
)
```

- Create adder function and load into csi
- Run on sci adder (+ 1 2 3 4)
- Run on sci (eval (+ 1 2 3 4))

We want: (+ (3 7 10 2))

```
(define (adder a_list)
  (cond
  ((null? a_list) 0)
  (else (eval(cons '+ a_list)))
)
```

Examples:

(adder '(1 2 3))

• We want: (+ (3 7 10 2))

Let's each write another way of doing this...

Create adder2 function and load into csi

Run on sci (adder2 '(3 7 10 2))

We want: (+ (3 7 10 2))

Let's each write another way of doing this... Hint: use car and cdr

Create adder2 function and load into csi

Run on sci (adder2 '(3 7 10 2))

Other functional languages

- Combination of many features of popular dialects of LISP, early 1980s
- Large and complex language, opposite of Scheme
- Features include: records; arrays; complex numbers; character strings; iterative control statements; etc.
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- Strongly typed (whereas Scheme is essentially typesless) with **no type coercions**

What were those?

- Syntax closer to Pascal and other imperative than to LISP
- Strongly typed (whereas Scheme is essentially typesless) with **no type coercions**
- Has identifiers, but once set cannot be changed more like final declarations in Java or const in C/C++

• Format:

fun name (parameters) = body;

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Example (run it):

fun circumf(r) = 3.14*r*r;

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fun circumf(r) = 3.14*r*r;

The type here is inferred as **float** from the type of the literal in the expression

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Example:

fun times $10(x) = 10^*x;$

Inferred as **int**

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Example:

fun square(x) = x^*x ;

Also inferred as **int** (default type)

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Example:

fun square(x) = x*x;

Also inferred as **int** (default type)

What happens if called with square(2.75)???

https://www.tutorialspoint.com/execute_smlnj_online.php

ML Language

Try running some code:

fun times10(x) = 10*x;
times10(5);

ML Language

Try running some code:

```
fun times10(x) = 10^*x;
times10(5);
```

times10(5.1);

What happens???

ML Language

Try running some code:

```
fun times10(x) = 10*x;
times10(5);
```

```
times10(5.1);
```

Yields error; expecting int...

It's strongly typed!!
ML Language

We could also specify type.

fun square(x:real) = x * x;

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We could also specify type.

fun square(x:real) = x * x;

Enough to infer that type is real

• These are all valid:

fun square(x:real) = x * x;

fun square(x) = (x:real) * x;

fun square(x) = x * (x:real)

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• These are all valid:

fun square(x:real) = x * x;

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fun square(x) = x * (x:real)

Enough to infer that type is real

Type inference also used in Haskell, Miranda, F#

ML Language

Try running some more code:

```
fun square(y:real) = y*y;
square(5.1);
square(5.0);
```

• What about this?

fun square(y:real) = y*y;
square(5);

• What about this?

```
fun square(y:real) = y*y;
square(5);
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```
Oops another type error...
```

• What about this?

```
fun square(y:real) = y*y;
square(5);
```

Oops another type error...

Note: user defined overloaded functions not allowed, so if we wanted a square function, one for real and one for int, would have to use different names...

• if else format:

if expression then expression

else else_expression

• Example:

fun fact (n:int) =

if n<=1 then 1
else n * fact(n-1);</pre>

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fun fact (n:int) =

if n<=1 then 1
else n * fact(n-1);</pre>

fact(4);

Run it...

Another way: pattern matching!

fun fact(0) = 1

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fun fact(0) = 1

| fact(1) = 1

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```
fun fact(0) = 1
```

```
| fact(1) = 1
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```
| fact(n:int) = n*fact(n-1);
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```
fun fact(0) = 1
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| fact(1) = 1
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| fact(n:int) = n*fact(n-1);
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Meant to mimic conditional functional definitions in math...

Another way: pattern matching!

```
fun fact(0) = 1
```

```
| fact(1) = 1
```

```
| fact(n:int) = n*fact(n-1);
```

Meant to mimic conditional functional definitions in math...

If param is int that is not 0 or 1 then third definition is used...

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Another way: pattern matching!

```
fun fact(0) = 1
```

```
| fact(1) = 1
```

```
| fact(n:int) = n*fact(n-1);
```

Note that don't need the **int** here since it is the default

Another way: pattern matching!

```
fun fact(0) = 1
```

```
| fact(1) = 1
```

```
| fact(n) = n*fact(n-1);
```

So this is also OK

Another way: pattern matching!

```
fun fact(0) = 1
| fact(1) = 1
| fact(n) = n*fact(n-1);
fact(4)
```

Let's try running code above...

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- Literal lists in brackets [3,5,7]; [] empty list
- :: used for cons
- 4::[3,5,7] evaluates to?
- [4,3,5,7]

- Try running these:
- 4::[3,5,7]
- hd([4,3,5,7])
- tl([4,3,5,7])

Number of elements in a list

fun length([]) = 0

Number of elements in a list

fun length([]) = 0

| length(h::t) = 1 + length(t);

Number of elements in a list

fun length([]) = 0

| length(h::t) = 1 + length(t);

length([1,3,5])

Try running it

fun append ([],lis2) = ?

(what should we write here?)

fun append ([],lis2) = lis2

fun append ([],lis2) = lis2

| append(h::t,lis2) = ?

What should we do?

fun append ([],lis2) = lis2

| append(h::t,lis2) = h::**?**

fun append ([],lis2) = lis2

| append(h::t,lis2) = h::append(t,lis2);

fun append ([],lis2) = lis2

| append(h::t,lis2) = h::append(t,lis2);

```
append([1,2],[3,4]);
```

Try running it...

Let's remind ourselves Scheme

```
(define (append lis1 lis2)
(cond
 ((null? lis1) lis2)
  (else (cons (car lis1)
      (append (cdr lis1) lis2)))
))
```

Reminding ourselves of cons (run it on csi):

(cons `(a b) `(c d)) returns ((a b) c d)

(cons `((a b) c) `(d (e f))) returns (((a b) c) d (e f))

```
fun append ([],lis2) = lis2
| append(h::t,lis2) =
h::append(t,lis2);
```

```
(define (append lis1 lis2)
(cond
 ((null? lis1) lis2)
  (else (cons (car lis1)
      (append (cdr lis1) lis2)))
))
```

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Let's each try fun adder

adder([1,2,3]) should return 6

Let's each try fun adder

fun adder([]) = 0

| adder (h::t)=h+adder(t);

adder([1,2,3,4,5]);

Names bound to values (constants)

• Format:

val new_name = expression;

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val new_name = expression;

Binds the value to name once and cannot be rebound (nothing like an assignment statement in an imperative language!)

Names bound to values (constants)

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val new_name = expression;

Example: usually used with a let statement:

```
fun area(radius) =
let val radius = 2.7
    val pi = 3.14159
in pi*radius*radius
end;
```

map

map(fn x =>x*x*x)[1,3,5];

map

map(fn x =>x*x*x)[1,3,5];

Note: different interpreters have slightly different notation; book notation different

Composing two functions

h = f o g

(lower case o)

Composing two functions

h = f o g

```
Example: (run it)
```

```
fun times10(x) = 10*x;
times10(5);
fun plus3(y) = 3 + y;
plus3(4);
val h = times10 o plus3;
h(7)
```