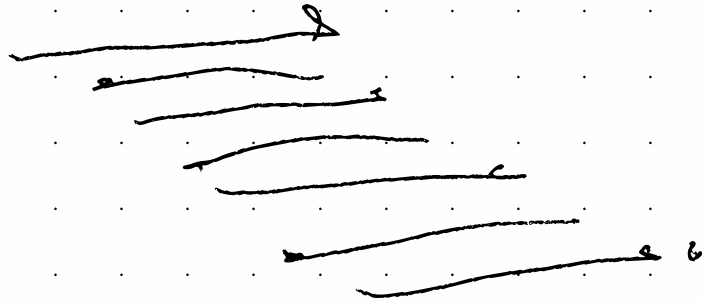


$O(n^2)$ ady

$n=10$

1 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1

$n/2$ }



$O(n^2)$

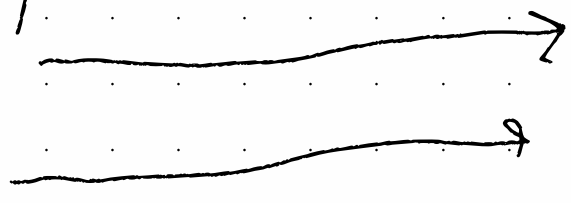
n

what does this mean

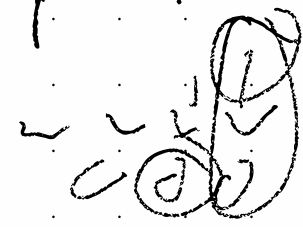
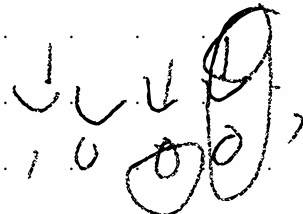
$O(n \log n)$

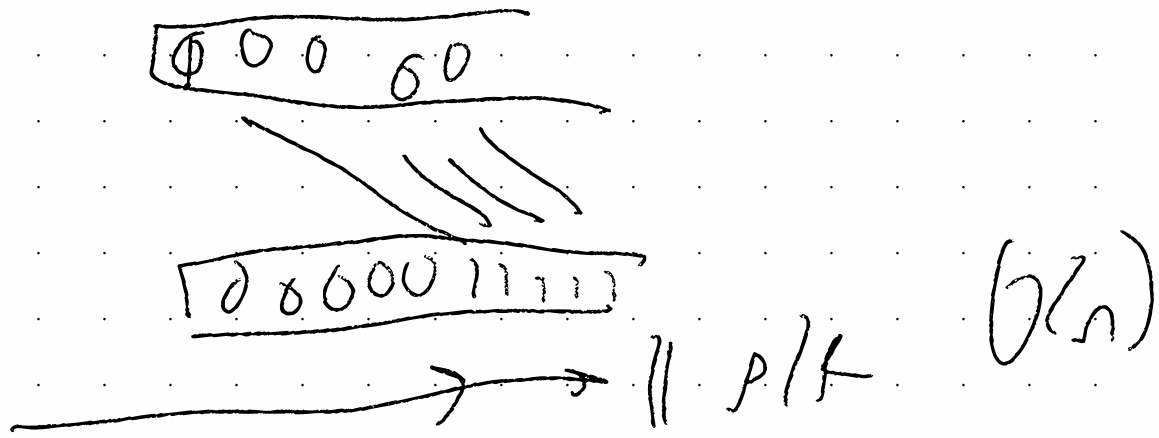
$\uparrow \phi \alpha \beta \phi \phi \gamma \gamma \gamma \gamma \gamma$

$\log n$



parity #0's
= parity #1's





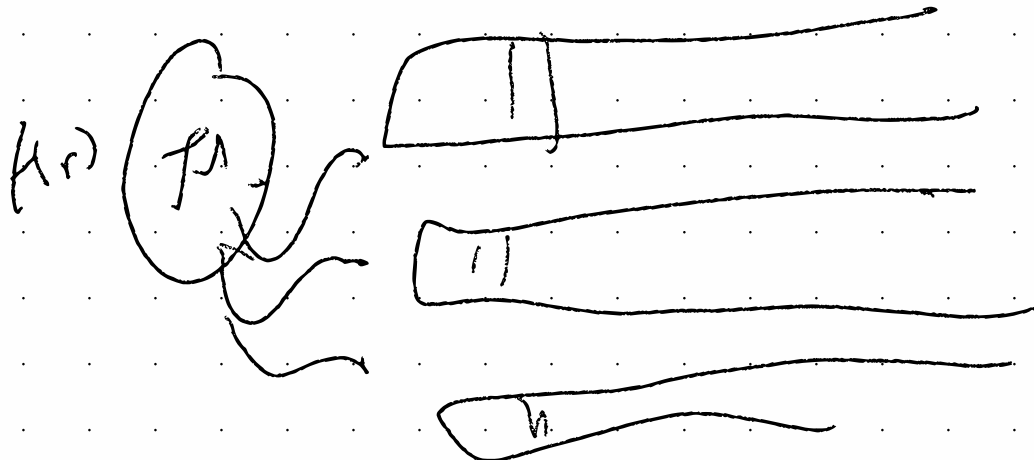
1 type $O(n^2)$, $O(n \log n)$

2 type $O(n)$

So how fast is recognition of

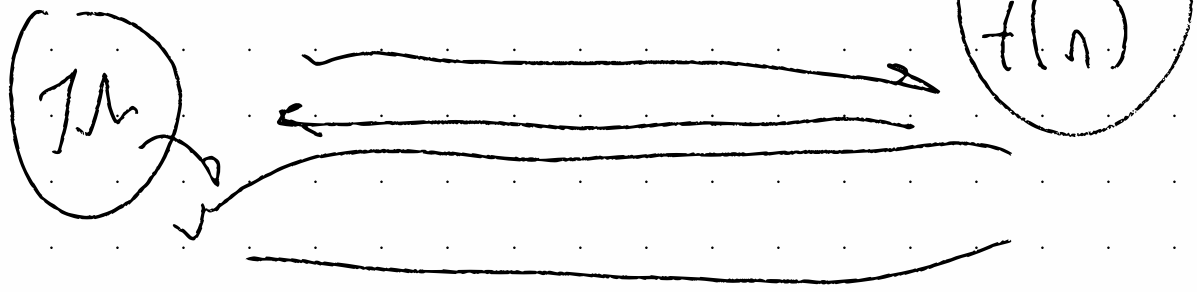
01^n ?

If a k -tape TM computes
in time $T(n)$, then a 1-tape
machine can compute in time $(T(n))^2$



over estimate

$f(n), f(n)$



all reasonable alt.

models of T^m , diff only

by a polynomial difference

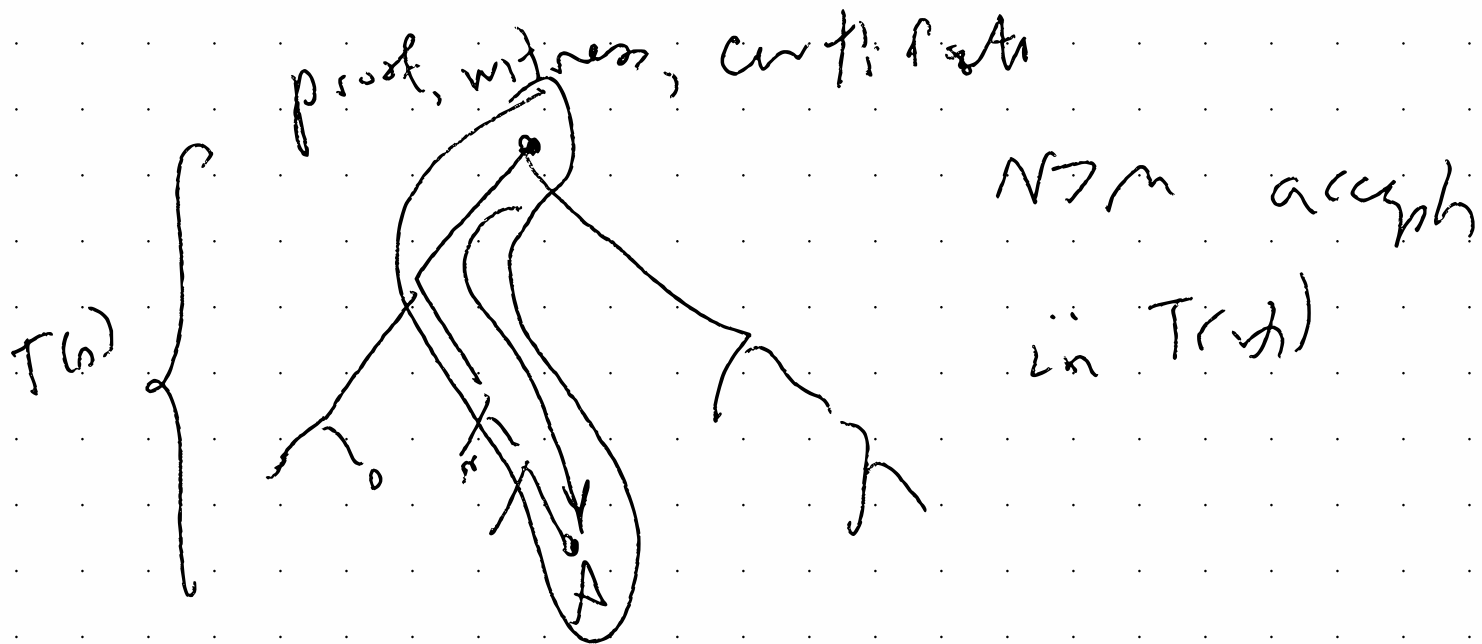
$P =$ all alg polynomial time
decidable

NP.

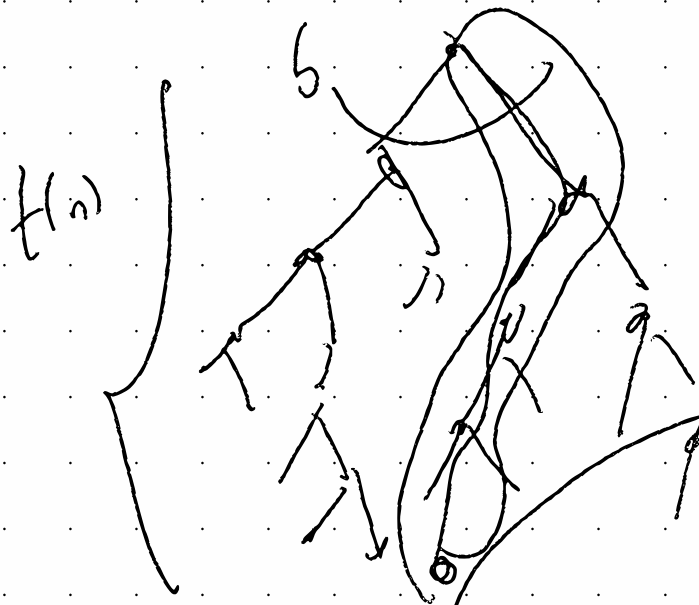
problem which an answer
can be verified by a P time
TM.

A NTM can be simulated by
a TM ~~is~~ in time $O(2^{T(n)})$

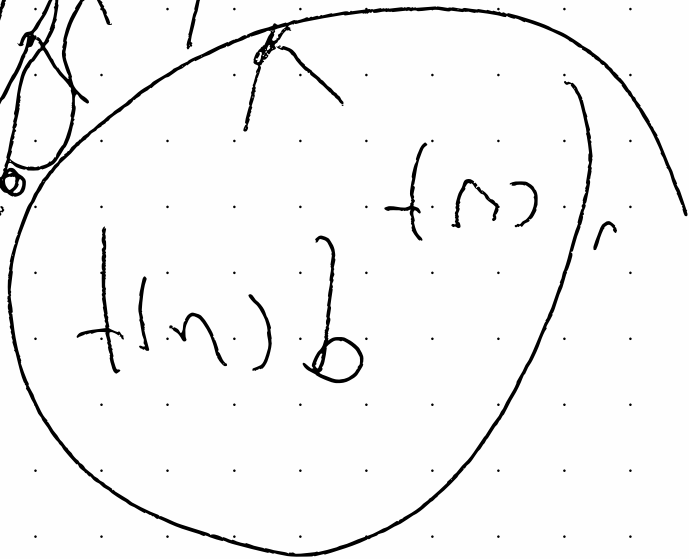
when NTM runs in time $T(n)$.



$\forall x \in L, \text{ATM}(x) = A$ is true
 $O(\sqrt{T(n)})$



↓ breadth first search



$$O(2^{f(n)})$$

