

an undecidable language

$$H(i, j) \equiv M_i(j)$$

is recognizable (a.k.a. recursively ~~known~~
enumerable)

$U(i, j) \leftarrow$ Universal Turing
machine

(U_i, \cdot) When Accept
provides a proof of
 y in the set
proof is the
computation

$\left\{ \begin{array}{l} \longrightarrow \text{reject (proof of } \notin) \\ \longrightarrow \text{ } \end{array} \right.$

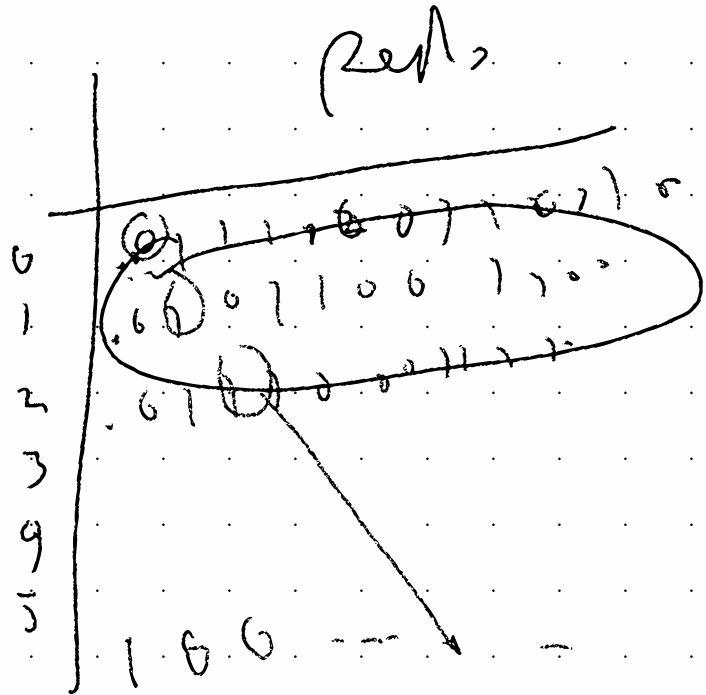
$\mathbb{R} \setminus \mathbb{H}$	1	2	3	3	
m_1	\textcircled{T}	\cancel{TF}	T	I	\cancel{TF}
m_2	T	$\textcircled{\cancel{TF}}$	I	\cancel{TF}	\cancel{TF}
m_3	I	\cancel{TF}	\textcircled{T}	\cancel{TF}	\cancel{TF}
\vdots	T	I	T	\textcircled{T}	\cancel{TF}
m_j	F	T	F	F	\dots

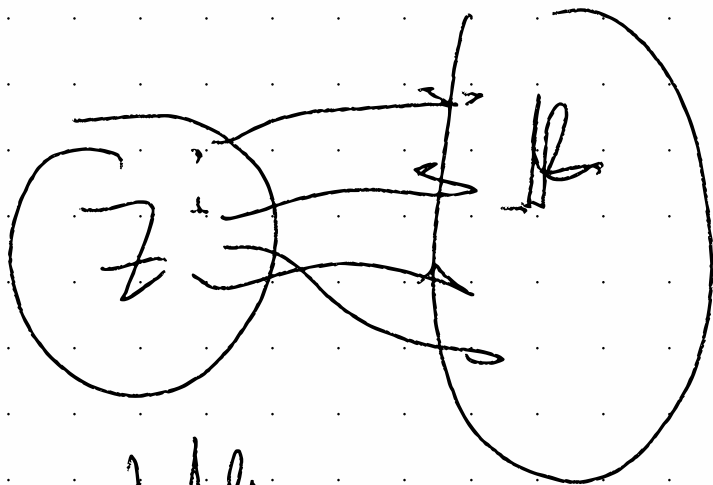
Diagonal algorithm program

how many real numbers are there?

$\mathbb{Q} \leftarrow$ infinite

\mathbb{Q}, \mathbb{N}





Un Countable

Countable

Q rational numbers
countable

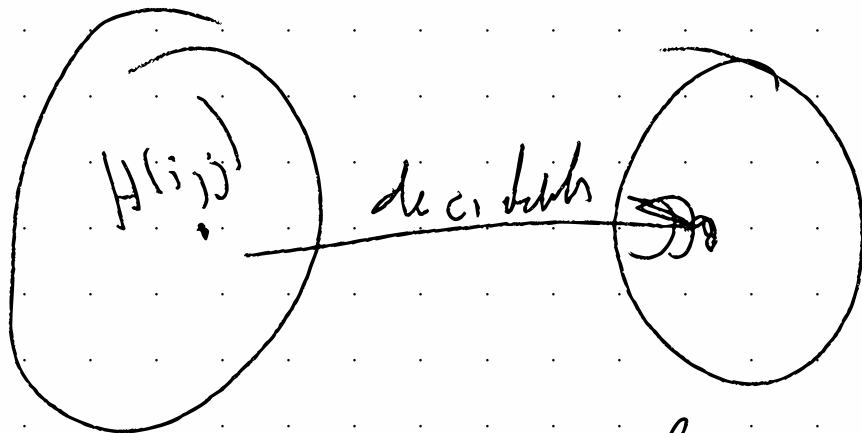
The set of \mathbb{N}^n is a countable set

0 1 2 3 ...
- 0 1 2 3 4 5 6 7 8 9 ...

{ } { 2 3 } { 7 } { 6 }

~~set~~ set of subsets of the integers

uncountable



\mathcal{A} instance of a
problem

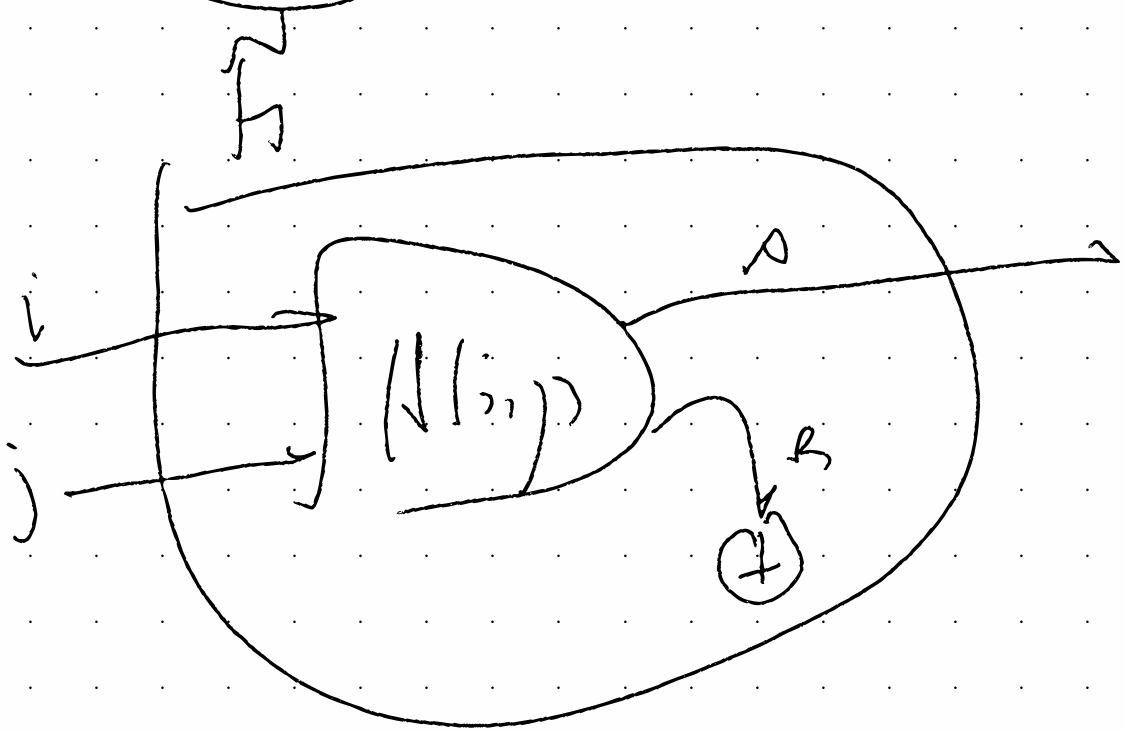
\mathcal{B} HALT(i, i)

$H(i, i)$

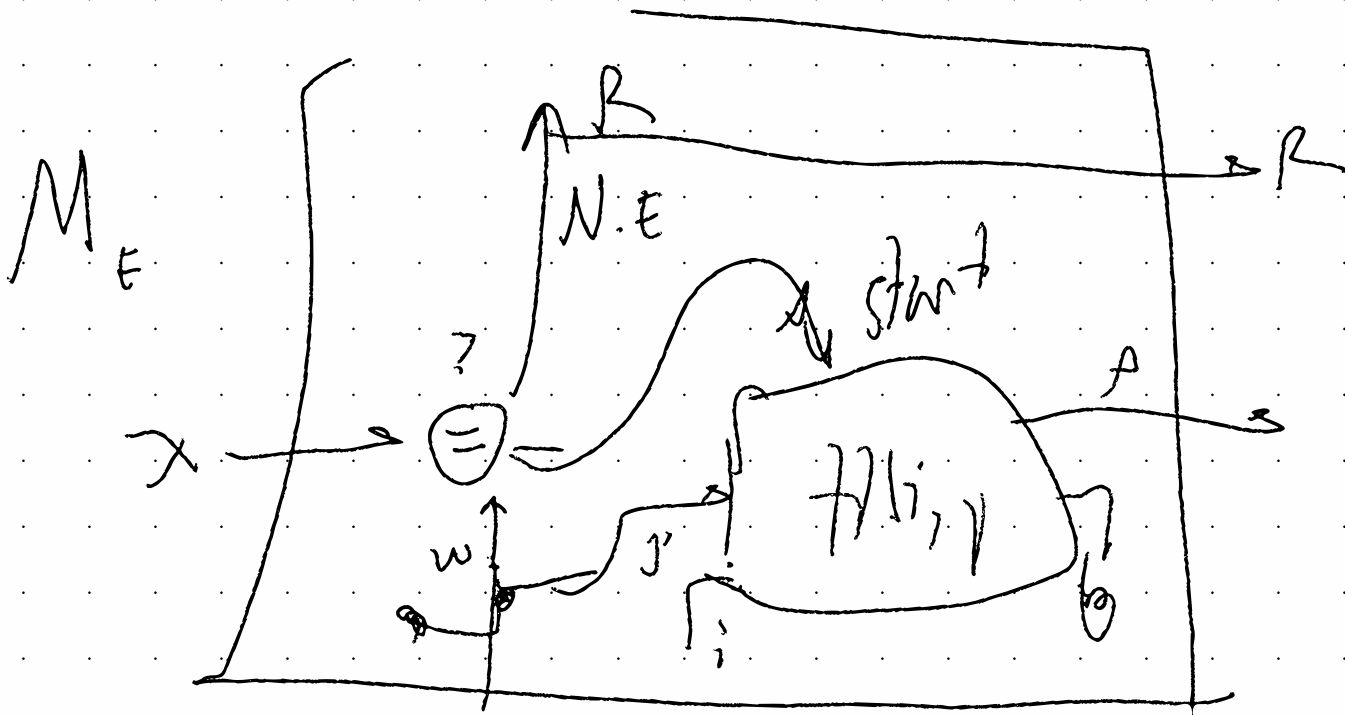
$= \text{HALT}(\hat{i}, \hat{i})$

HALT

is undecidable



$$L(M_\epsilon) \neq \emptyset \iff h(\cdot, \cdot) = A$$



$\text{if } (A, i, j) = \Delta \longrightarrow \Sigma^* \text{ regular}$
 $\neq \Delta \longrightarrow \text{ok, not regular}$

