#### Data Structures and Algorithm Analysis (CSC317)

#### **Dynamic Programming 1**

**Odelia Schwartz** 

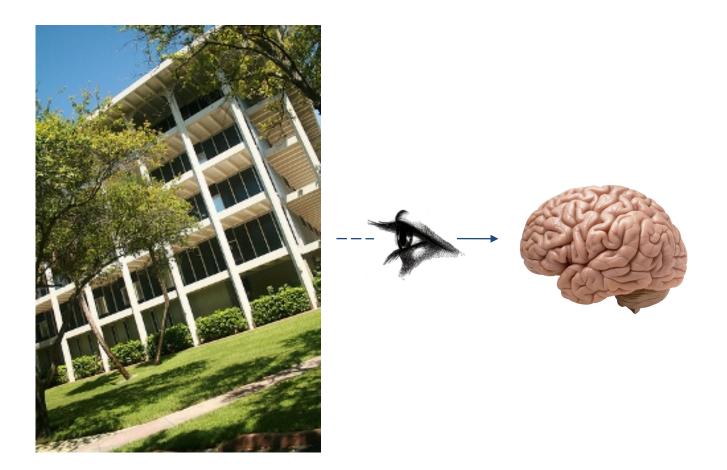
#### CSC317 House Keeping

• Introductions...

Your major and what do you hope to get out of this course?

#### In my field... Computational neuroscience

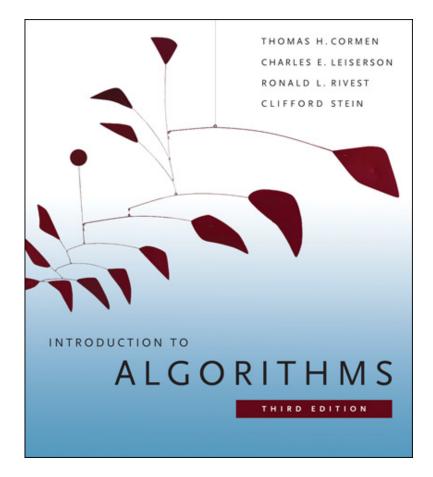
Brain receives input, processes information, and computes outputs. What algorithms does the brain use??



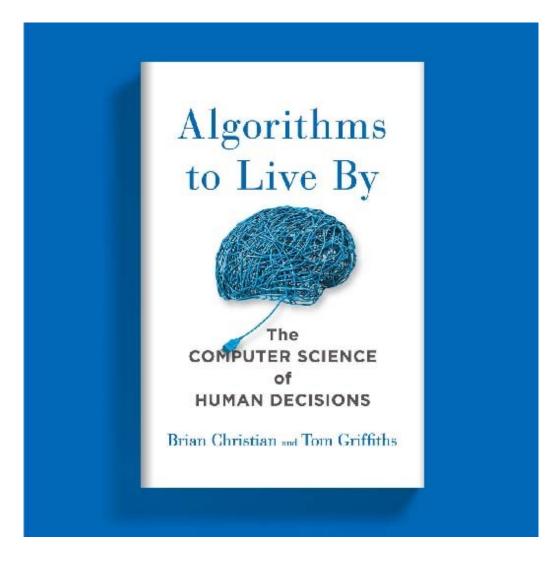
#### CSC317 House Keeping

- Course homepage: I will post slides:
   <u>http://www.cs.miami.edu/home/odelia/teaching/csc317\_fall19/index.html</u>
- My typed slides will be posted on a regular basis; in class develop on the board...
- Odelia Schwartz (odelia at cs miami dot edu). Encouraged to email to make appointment or stop by
- Assignments continue to be on BB
- Continued structure of quizzes (and no final!)

# Data Structures and Algorithm Analysis (CSC317)



#### Optional reading, beyond scope



• Divide and Conquer

• Divide and Conquer

Next:

- Dynamic Programming
- Greedy algorithms

• Divide and Conquer

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- Greedy algorithms

# **Dynamic Programming**

- General, powerful
- Problems that may naively have exponential running time, but can be made poynomial (fast!)
- "Programming"?

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- "Programming"? Not programming languages; Bellman was interested in planning and decision making. See:

http://www.cs.miami.edu/home/odelia/teaching/csc317\_fall19/syllabus/dy\_birth.pdf

# **Dynamic Programming**

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 Can be thought of as "tabular programming" as in "table." Main approach: hold answers to previous problems already solved in a table, so that can be used again without recomputing.

# Dynamic Programming Class Outline

- Examples of applications (motivation)
- Simple example to gain intuition
- Back to applications and more examples (next classes)

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- Computational Biology (genome similarity; also spike similarity; file text similarity)
- Cypher to Thomas Jefferson (will mention)
- Finding shortest path (later)

• Computational Biology (genome similarity)

Strings from alphabet {A, C, G, T}

Example: ACGGAT CCGCTT

How to determine similarity?

And why? Understanding one genome sequence and its similarity to another can teach us about function...

• Computational Biology (genome similarity)

Strings from alphabet {A, C, G, T}

Example: ACGGAT CCGCTT

How to determine similarity?

Number of changes from one to another small Allowed to change character Find the Longest Common Subsequence

• Computational Biology (genome similarity)

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What is the Longest Common Subsequence?

• Computational Biology (genome similarity)

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Example: ACGGAT CCGCTT

What is the Longest Common Subsequence?

Answer: 3 CGT

Is answer unique?

• Computational Biology (genome similarity)

Strings from alphabet {A, C, G, T}

Example: ACGGAT CCGCTT

What is the Longest Common Subsequence?

We easily eye balled answer for these short sequences. Longer sequences of 500 or more characters? Brute force solution?

• Computational Biology (genome similarity)

What is the Longest Common Subsequence? ACCCGGTCGAGTG... GTCGTTCGGAATT...

Brute force: Try all subsequences in 1<sup>st</sup> string and compare to second string... n=500 then 2^500 possibilities

Pick first character or do not...
Pick 2<sup>nd</sup> character or do not...
Pick 3<sup>rd</sup> or do not...
2 \* 2 \* 2 \* 2 .... \* 2 (n times)

• Computational Biology (genome similarity)

What is the Longest Common Subsequence? ACCCGGTCGAGTG... GTCGTTCGGAATT...

Brute force: n=500 then 2^500 possibilities

Pick  $1^{st}$  character or do not... Pick  $2^{nd}$  character or do not... Pick  $3^{rd}$  or do not... 2 \* 2 \* 2 \* 2 ... \* 2 (n times) =  $2^{n}$ Actually need to multiply by length of  $2^{nd}$  string

• Computational Biology (genome similarity)

What is the Longest Common Subsequence? ACGGAT CCGCTT

We learned Divide and Conquer. Will this approach work?

• Computational Biology (genome similarity)

What is the Longest Common Subsequence? A C G G A T C C G C T T

We learned Divide and Conquer. Will this approach work?

Answer: No. Not in a simple way.

It could for this example, but not generally...

- ACG GAT
- CCG CTT
  - CG T

• Computational Biology (genome similarity)

What is the Longest Common Subsequence? CGTGAC CGGTTT

We learned Divide and Conquer. Will this approach work?

Answer: No. Not in a simple way. Does not find C G T, Unless you look across the midline...

Doesn't work easily here...

- CGT GAC
- CGG TTT

• Computational Biology (genome similarity)

What is the Longest Common Subsequence? CGTGAC CGGTTT

We will learn a Dynamic Programming approach...

• Spike Similarity...

Cypher to Thomas Jefferson

http://www.cs.miami.edu/home/odelia/teaching/ csc317\_fall19/syllabus/cipherJefferson-amsci2009-03S.pdf

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Figure 1. On December 19, 1807, Robert Pat terson (*uir left*)—a professor of mathematic at the University of Permsylvania—mote letter to Thomas Jefferson (*unmediati left*) Patterson described his vision of a "prefect cipher," which required four elements adaptable to all Inauguages, easy to memorize simple to perform and inscrutable withou the key. Fatterson also described an encryp tion technique that he believed met these criteria. In addition, he included encryptet text, which he said could never be decrypted text, which he said out neve the decryptet text, which he said out never be decrypted to be charged and never be decrypted to son's challenge using techniques that could have been applied—if laboriously—in the any full neutry. (All letter reproduction

• Cypher to Thomas Jefferson

http://www.cs.miami.edu/home/odelia/teaching/ csc317\_fall19/syllabus/cipherJefferson-amsci2009-03S.pdf

Cypher to Thomas Jefferson

http://www.cs.miami.edu/home/odelia/teaching/ csc317\_fall19/syllabus/cipherJefferson-amsci2009-03S.pdf

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Cypher to Thomas Jefferson ullet

http://www.cs.miami.edu/home/odelia/teaching/ csc317\_fall19/syllabus/cipherJefferson-amsci2009-03S.pdf

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5	<mark>a</mark> psevh	83	stlrcwreh
6	penwee	14	sees <mark>b</mark> inlei
7	aaoobc	62	arpenwee
8	rcwreh	20	uvclst

# Simple example (to build intuition)

• Fibonacci!

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Fib(n)

- 1. If n==0 return 1
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- 3. else return Fib(n-1) + Fib(n-2)

Good algorithm??

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• Fibonacci!

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Good algorithm?? Does the job but ... no, very wasteful! Why?

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A lot of recomputing ...

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Good algorithm?? Does the job but ... no, very wasteful! Why?

Keep repeating computations ... Fib(25) = Fib(24) + Fib(23) ... Fib(24) = Fib(23) + Fib(22)...

Fib(n)

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Recursion tree on the board...

Fib(n)

- 1. If n==0 return 1
- 2. If n==1 return 1
- 3. else return Fib(n-1) + Fib(n-2)

See animation:

https://www.cs.usfca.edu/~galles/visualization/DPFib.html

# Simple example (Fibonacci)

Summary so far:

- Overlapping subproblems (lots)!
- Solution to big problem can be constructed from solutions to subproblems
- Example of type of problems that can be solved with Dynamic Programming

# Simple example (Fibonacci)

Dynamic Programming Fibonacci:

- Main idea: Save in dictionary (e.g., array) subproblems already solved, so no need to recompute
- Memoization: from memo pad or memory ... funky name...

On the board...

a. Initialization: Let mem be a new array with values initialized to minus infinity

- a. Initialization: Let mem be a new array with values initialized to minus infinity
- b. Fib(n) //Memoized Dynamic Programming
  - 1. If mem[n]>=0
  - 2. return mem[n] //if already previously computed in memo pad
  - 3. if n==0 return 1
  - 4. if n==1 return 1
  - 5. else f = Fib(n-1) + Fib(n-2) //otherwise compute and save value
  - 6. mem[n] = f //save value in memo pad
  - 7. return f

Plot tree: On the board...

- Run time proportional to n
- Second time encountered, just use memoized result...
- Cuts off whole recursion subtrees!

# of subproblems: n (size of array) work per subproblem: constant

Plot tree: On the board...

- Run time proportional to n
- Second time encountered, just use memoized result...
- Cuts off whole recursion subtrees!

See online by Galles:

https://www.cs.usfca.edu/~galles/visualization/DPFib.html

See online by Rosenberg:

http://www.cs.miami.edu/home/odelia/teaching/fib2019.html

Summary: Recursion + memoization

Fib(n) //Bottom-up Dynamic Programming 1. Let mem[0..n] be a new array 2. mem[0] = 1 3. mem[1] = 1 4. For i=2 to n 5. mem[i] = mem[i-1] + mem[i-2]

6. return mem[n]

Question: Is bottom-up algorithm the same or different from the previous recursive memoized solution?

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Answer: One for loop, complexity proportional to n Equivalent solution to recursive memoization (same things Happen in same order; but in bottom-up we know and set the order in advance)

Question: If this is how we were first taught Fibbonacci, why bother with naïve inefficient recursion, memoized solution, etc. first?

Answer: Other problems initially less intuitive, but approach will be similar (think back to Genome question)

### Dynamic Programming so far

- 1. Overlapping subproblems (same subproblems solved over and over again
- 2. Solution to big problem constructed from solutions to smaller subproblems (optimal substructure; more on later)
- To make algorithm more efficient, we either memoized (saved solutions to smaller subproblems in a table) as we recursed; or we saved solutions to subproblems bottom-up. These turned out equivalent.

### Dynamic Programming so far

Question: Both Dynamic Programming and Divide & Conquer have recursive solutions. But they are different. Why?

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Question: Both Dynamic Programming and Divide & Conquer have recursive solutions. But they are different. Why?

Answer: For instance, Divide & Conquer doesn't have overlapping subproblems...

### Next

- In Fib clear what the smaller subproblems are, and how knowing their solution solves the bigger problem
- Start to build intuition with more complex problems, starting from genome similarity and Longest Common Subsequence...